

## Optimal spliced alignment of homologous cDNA to a genomic DNA template

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### Abstract

**Motivation:** Supplementary cDNA or EST evidence is often decisive for discriminating between alternative gene predictions derived from computational sequence inspection by any of a number of requisite programs. Without additional experimental effort, this approach must rely on the occurrence of cognate ESTs for the gene under consideration in available, generally incomplete, EST collections for the given species. In some cases, particular exon assignments can be supported by sequence matching even if the cDNA or EST is produced from non-cognate genomic DNA, including different loci of a gene family or homologous loci from different species. However, marginally significant sequence matching alone can also be misleading. We sought to develop an algorithm that would simultaneously score for predicted intrinsic splice site strength and sequence matching between the genomic DNA template and a related cDNA or EST. In this case, weakly predicted splice sites may be chosen for the optimal scoring spliced alignment on the basis of surrounding sequence matching. Strongly predicted splice sites will enter the optimal spliced alignment even without strong sequence matching.

**Results:** We designed a novel algorithm that produces the optimal spliced alignment of a genomic DNA with a cDNA or EST based on scoring for both sequence matching and intrinsic splice site strength. By example, we demonstrate that this combined approach appears to improve gene prediction accuracy compared with current methods that rely only on either search by content and signal or on sequence similarity.

**Availability:** The algorithm is available as a C subroutine and is implemented in the SplicePredictor and GeneSequer programs. The source code is available via anonymous ftp from ftp.zmdb.iastate.edu. Both programs are also implemented as a Web service at <http://gremlin1.zool.iastate.edu/cgi-bin/sp.cgi> and <http://gremlin1.zool.iastate.edu/cgi-bin/gs.cgi>, respectively.

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### Introduction

Global sequencing efforts are currently producing vast amounts of raw genomic sequence data for many different organisms. The pace of sequencing necessitates that the sequence annotation, in particular with respect to gene structure, be largely based on computational algorithms for automated sequence interpretation [for a recent review see Claverie (1997)]. Experimental evidence for exon assignments may derive from cDNA or EST sequencing. Typically, the cDNA sequences will come from independently sequenced cDNA libraries, and assignment of a cDNA to its cognate gene will be on the basis of sequence identity. In the simplest, unambiguous case, the alignment will consist of (1) perfectly matching segments corresponding to the exons, and (2) deletions in the cDNA corresponding to introns in the genomic template. In practice, matching may be less than perfect due to either sequencing errors or, more importantly, due to matching of genomic sequences with non-cognate cDNA. The non-cognate cDNAs derive not from the given genetic locus but from homologous loci, for example, the corresponding locus in a related species or a duplicated locus representing a different member of the same gene family. In this case, the alignment will generally have to include mismatches and gaps, but may still strongly support a particular gene structure prediction at the locus being analyzed.

We present the subroutine sahmtD (Spliced Alignment Hidden Markov Tool for cDNA) which implements a dynamic programming algorithm to efficiently calculate the optimal scoring alignment between an assumed template DNA and a second sequence representing a related collinear spliced product. The novelty in our approach compared to previous algorithms (Gotoh, 1982; Florea *et al.*, 1998; Huang, 1994; Huang *et al.*, 1997; Mott, 1997) consists (1) in the simultaneous assessment of the significance of the sequence alignment and the intrinsic

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quality of the implied splice sites, and (2) in the explicit assignment of exon or intron status to each nucleotide in the genomic DNA. The algorithm is considerably more reliable in cases where global sequence similarity is weak or compromised by regions of poor local similarity. Applications are illustrated in the context of resolution of multiple hits in cDNA database searches with genomic sequence queries and the study of a hypothetical novel *Arabidopsis thaliana* gene family.

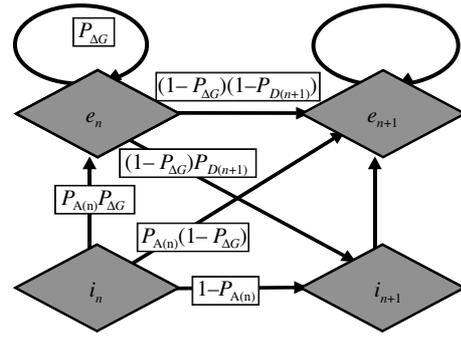
### System and methods

We pose the problem of finding an optimal alignment of a genomic nucleotide sequence  $G_1, G_2, \dots, G_N$  of length  $N$  with a cDNA or EST nucleotide sequence  $C_1, C_2, \dots, C_M$  of length  $M$ . A precise definition will be given later of optimality relative to a scoring system that simultaneously evaluates the pairwise sequence similarity and the quality of predicted splice sites in the genomic sequence. Both sequences consist of letters from the alphabet  $\mathbf{A} = \{A, C, G, T, N\}$  where A, C, G, T denote the nucleotides adenine, cytosine, guanine, and thymine, respectively, and N denotes an undetermined nucleotide. An alignment between the sequences may include gaps in either sequence, indicated by the additional gap symbol ‘\_’ juxtaposed to each of the letters comprising the corresponding insertion in the other sequence. We use the notation  $\mathbf{A}^+$  for the alphabet superset  $\{A, C, G, T, N, \_ \}$  and  $\mathbf{A}^-$  for the subset  $\{A, C, G, T\}$ . All possible alignments may be viewed as outputs of a Hidden Markov Model (HMM). The HMM defines a probability space consisting of all possible “threadings” of cDNA sequences of length  $M$  over the alphabet  $\mathbf{A}$  into the given genomic sequence. The formulation of the algorithm in terms of a HMM is merely for convenience of presentation. The coding of the algorithm involves log probabilities that are in practice replaced by any suitable additive weights without loss of generality.

The state sequence underlying a given alignment will be denoted as  $Q = q_1 q_2 \dots q_L$ , where  $\max\{M, N\} \leq L \leq M + N$ . The set of states of the HMM consists of the exon states  $e_n, n = 1, 2, \dots, N$ , with output  $\begin{matrix} X \\ Y \end{matrix}$ ,  $X, Y \in \mathbf{A}^+$ , and the intron states  $i_n, n = 1, 2, \dots, N$ , with output  $\begin{matrix} G_n \\ \star \end{matrix}$ , where the  $\star$  symbol serves as a placeholder for the spliced sequence parts. Transitions between the states are limited to the following transitions with non-zero transition probabilities  $\tau_{q_l, q_{l+1}}$  (Figure 1):

$$\begin{aligned} \tau_{e_n, e_{n+1}} &= (1 - P_{\Delta G})(1 - P_{D(n+1)}) \\ \tau_{i_n, e_{n+1}} &= P_{A(n)}(1 - P_{\Delta G}) \\ \tau_{e_n, e_n} &= P_{\Delta G} & \tau_{i_n, e_n} &= P_{A(n)}P_{\Delta G} \\ \tau_{e_n, i_{n+1}} &= (1 - P_{\Delta G})P_{D(n+1)} & \tau_{i_n, i_{n+1}} &= 1 - P_{A(n)} \end{aligned}$$

for  $n = 1, 2, \dots, N$  (third line) or  $N - 1$  (other



**Fig. 1.** States and transitions of the Hidden Markov Model. States are represented by diamonds. The model involves exon ( $e$ ) and intron ( $i$ ) states. The index  $n$  represents the position in the genomic sequence assigned to the state. Transitions between the states are indicated by arrows. The transition probabilities are shown for transitions from states  $e_n$  and  $i_n$ .  $P_{\Delta G}$  is the probability of a nucleotide deletion in the genomic sequence.  $P_{D(n)}$  and  $P_{A(n)}$  are the probabilities of position  $n$  in the genomic DNA to be a donor or acceptor site, respectively.

lines). Here  $P_{D(n)}$  and  $P_{A(n)}$  are the pre-determined probabilities that  $G_n$  in the genomic sequence is the first base (donor site) or last base (acceptor site) of an intron, respectively. In the applications for plant gene identification discussed here, these values are set equal to the  $P$ -values calculated by the SplicePredictor program (Brendel and Kleffe, 1998; Kleffe *et al.*, 1996). Sites that are not scored by SplicePredictor are given small positive probabilities so that non-consensus sites supported by surrounding exon sequence matching are not excluded *a priori*. Other assignments could be made, for example derived from NetPlantGene output (Hebsgaard *et al.*, 1996) or (in the absence of models appropriate for the given species) generic assignments (distinguishing only between GT, GC, and other potential donor sites, and between AG and other potential acceptor sites).  $P_{\Delta G}$  is a parameter that denotes the probability of inserting a gap symbol in the genomic sequence.

The output weights in the exon states  $e_n$  are set to

$$\begin{aligned} \log P_{e_n} \left( \begin{matrix} G_n \\ X \end{matrix} \right) &= \begin{cases} \sigma & \text{if } G_n = X \\ \mu & \text{otherwise} \end{cases} \\ \log P_{e_n} \left( \begin{matrix} N \\ X \end{matrix} \right) &= \nu \\ \log P_{e_n} \left( \begin{matrix} G_n \\ N \end{matrix} \right) &= \nu \\ \log P_{e_n} \left( \begin{matrix} N \\ N \end{matrix} \right) &= \nu \\ \log P_{e_n} \left( \begin{matrix} G_n \\ - \end{matrix} \right) &= \delta \end{aligned}$$

$$\log P_{e_n} \begin{pmatrix} N \\ - \end{pmatrix} = \delta$$

for  $G_n \in \mathbf{A}^-$ ,  $X \in \mathbf{A}^-$ , where  $\sigma$ ,  $\mu$ ,  $\nu$ , and  $\delta$  represent the weights for identities, mismatches, alignment positions involving undetermined characters, and cDNA deletions, respectively. The output weights corresponding to genomic sequence deletions are also set uniformly to

$$\log P_{e_n} \begin{pmatrix} - \\ X \end{pmatrix} = \delta, \quad X \in \mathbf{A}.$$

Note that for a strict HMM formulation genomic sequence deletions would be output from additional ‘delete’ states. However, because the transitions from the delete states are exactly like the transitions from the corresponding exon states, our formulation is more efficient (in the coding detailed below, the output weights are always assigned in conjunction with the transition probabilities so that it is always clear whether  $e_n$  corresponds to a delete state or not). For the intron states  $i_n$ ,

$$\log P_{i_n} \begin{pmatrix} G_n \\ \star \end{pmatrix} = 0.$$

Thus the output probabilities involve only four parameters (see Implementation).

With the above formulation, optimal alignments are precisely defined as state sequences  $Q = q_1 q_2 \dots q_L$  with associated output  $S_M^N$  (representing a sequence alignment of  $G_1 G_2 \dots G_N$  with  $C_1 C_2 \dots C_M$ ) such that the joint probability  $P(Q, S_M^N)$  is maximal over all possible  $Q$  and  $S_M^N$ . This maximal probability is calculated in standard fashion as

$$P = \max\{E_M^N, I_M^N\},$$

where

$$E_m^n = \max P(Q = q_1 q_2 \dots q_l, q_l = e_n, S_m^n),$$

and

$$I_m^n = \max P(Q = q_1 q_2 \dots q_l, q_l = i_n, S_m^n),$$

for  $n = 1, 2, \dots, N$ ,  $m = 1, 2, \dots, M$ ,  $\max\{m, n\} \leq l \leq m + n$ , and maximization is over all possible  $Q$  and  $S_m^n$  representing alignments of  $G_1, G_2, \dots, G_n$  with  $C_1, C_2, \dots, C_m$ .

Let  $\tau_{e_0, e_1} = \tau_{i_0, e_1} = \tau_{e_1}$  and  $\tau_{e_0, i_1} = \tau_{i_0, i_1} = \tau_{i_1} = 1 - \tau_{e_1}$ , where  $\tau_{e_1}$  is the initial exon state probability. Then  $E_M^N$  and  $I_M^N$  are found from the following recursion:

$$\begin{aligned} E_0^n &= I_0^n = 1, \\ E_m^0 &= 1, \quad I_m^0 = 0, \quad n = 0, 1, \dots, N, m = 1, 2, \dots, M, \\ E_m^n &= \max \left\{ \max \left\{ E_{m-1}^{n-1} \tau_{e_{n-1}, e_n}, I_{m-1}^{n-1} \tau_{i_{n-1}, e_n} \right\} P_{e_n} \begin{pmatrix} G_n \\ - \end{pmatrix}, \right. \\ &\quad \left. \max \left\{ E_{m-1}^n \tau_{e_n, e_n}, I_{m-1}^n \tau_{i_n, e_n} \right\} P_{e_n} \begin{pmatrix} - \\ C_m \end{pmatrix} \right\}, \end{aligned}$$

$$\begin{aligned} &\max \left\{ E_{m-1}^{n-1} \tau_{e_{n-1}, e_n}, I_{m-1}^{n-1} \tau_{i_{n-1}, e_n} \right\} P_{e_n} \begin{pmatrix} G_n \\ C_m \end{pmatrix}, \\ &\max \left\{ E_{m-1}^n \tau_{e_n, e_n}, I_{m-1}^n \tau_{i_n, e_n} \right\} P_{e_n} \begin{pmatrix} - \\ C_m \end{pmatrix} \right\}, \end{aligned}$$

$$\begin{aligned} I_m^n &= \max \left\{ E_{m-1}^{n-1} \tau_{e_{n-1}, i_n}, I_{m-1}^{n-1} \tau_{i_{n-1}, i_n} \right\}, \\ &n = 1, 2, \dots, N, m = 1, 2, \dots, M. \end{aligned}$$

At each maximization step, the state transition and output yielding the maximum are stored to facilitate the backtracing of the optimal alignment(s). See Figure 2 for a hypothetical example.

### Implementation

Given the  $P$ -values according to the SplicePredictor, few parameters need to be specified for a complete implementation of the algorithm. The following default values worked well for a large range of applications we examined:

$$\begin{aligned} \tau_{e_1} &= 0.5 \\ P_{\Delta G} &= 0.03 \\ \sigma &= 2.0 \quad \mu = -2.0 \quad \nu = 0.0 \quad \delta = -4.0. \end{aligned}$$

In a typical application the genomic DNA template will extend 5' and 3' of the cDNA ends. Setting  $E_0^n$  and  $I_0^n$  to one for all  $n$  amounts to no end gap penalties at the 5' end. To symmetrize with respect to 3' end gap penalties, in the programmed updating of  $E_m^n$  the output weight

$$\log P_{e_n} \begin{pmatrix} G_n \\ - \end{pmatrix}$$

is set to zero for  $m = M$ . Similarly,

$$\log P_{e_n} \begin{pmatrix} - \\ C_m \end{pmatrix}$$

is set to zero for  $n = N$ . Note that within-exon gaps in the genomic DNA are more costly than gaps in the cDNA because  $\log P_{\Delta G} < \log(1 - P_{\Delta G})$ , a desirable setting as cDNA or EST sequences are typically less reliably determined in practice.

The memory requirements of the program are minimal because the updating of the  $E_m^n$  and  $I_m^n$  matrices at a given index pair only involves cells at most one row and column up and to the left. In practice, for given index  $n$  the program simply fills out one of two row arrays of size  $M$  (labeled  $n \bmod 2$ ), using information from the previous calculations stored in the array labeled  $(n - 1) \bmod 2$ . In the next step the then unnecessary information in array  $(n - 1) \bmod 2$  is overwritten. For convenience, we store the maximal scoring state transitions for backtracing an optimal path. This allows rapid recovery of an optimal alignment at the cost of extra storage.

<b>index <math>n</math></b>	1	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
<b>genomic DNA</b>	T	-	C	A	G	G	T	A	A	G	T	C	A	A	A	T
<b>EST</b>	T	T	C	A	N	*	*	*	*	*	T	C	-	-	C	T
<b>index <math>m</math></b>	1	2	3	4	5	5	5	5	5	5	6	7	7	7	8	9
<b>state sequence</b>	$e_1$	$e_1$	$e_2$	$e_3$	$e_4$	$i_5$	$i_6$	$i_7$	$i_8$	$i_9$	$e_{10}$	$e_{11}$	$e_{12}$	$e_{13}$	$e_{14}$	$e_{15}$
<b>transition probabilities</b>	$\tau_{e1}$	$\tau_{e1e1}$	$\tau_{e1e2}$	$\tau_{e2e3}$	$\tau_{e3e4}$	$\tau_{e4i5}$	$\tau_{i5i6}$	$\tau_{i6i7}$	$\tau_{i7i8}$	$\tau_{i8i9}$	$\tau_{i9e10}$	$\tau_{e10e11}$	$\tau_{e11e12}$	$\tau_{e12e13}$	$\tau_{e13e14}$	$\tau_{e14e15}$
<b>output weights</b>	$\sigma$	$\delta$	$\sigma$	$\sigma$	$\nu$	0	0	0	0	0	$\sigma$	$\sigma$	$\delta$	$\delta$	$\mu$	$\sigma$

**Fig. 2.** Hypothetical alignment of a genomic DNA with an EST sequence. The genomic sequence (second row) comprises 15 nucleotides. The EST sequence (third row) is of length nine nucleotides, including in position 5 a non-determined base (N). An alignment is shown that assigns intron status to the genomic DNA positions 5–9. The underlying state sequence is displayed in the fifth row. Indices  $n$  (first row) and  $m$  (fourth row) record the position in the genomic DNA and EST sequences, respectively. A deletion in the genomic sequence is accommodated by the transition from state  $e_1$  into itself. The transition probabilities and output weights (bottom two rows) are assigned as described in the text. The algorithm maximizes the sum of log transition probabilities plus output weights over all possible spliced alignments.

The algorithm was implemented as the C subroutine sahmtD (for Spliced Alignment Hidden Markov Tool for cDNA) in our previous SplicePredictor program (Brendel and Kleffe, 1998; Kleffe *et al.*, 1996). Limitations on the maximal lengths of the genomic DNA and cDNA depend on the memory of the CPU. Our WWW server is currently set up to align genomic DNA segments up to 13 kb against a cDNA of up to 7 kb. If the input exceeds these limits, an exit warning is displayed and recompilation suggested with increased limits. For plant genes that typically lack long exons and introns (90% of maize and *Arabidopsis* exons and introns are less than 510 nucleotides; Brendel *et al.*, 1998) these limits appear adequate in practice. Detection of long introns would not be explicitly feasible by our approach. However, a long intron would be a reasonable interpretation if the 5' and 3' ends of a single EST matched dispersed regions in a genomic DNA.

One possible use of the algorithm is to screen a novel large genomic contig against an entire EST database. This could be achieved by pre-selecting matching ESTs with a fast screening program like BLAST (Altschul *et al.*, 1997) or its derivatives; e.g. <http://genome-www2.stanford.edu/cgi-bin/AtDB/nph-blast2atdb>. This strategy has been pursued by Florea *et al.* (1998) and Mott (1997). Alternatively, we have also implemented the sahmtD subroutine in a standalone program called GeneSeqer. In GeneSeqer, each EST is initially fast-screened against the genomic DNA for a matching region of specific quality. A matching region defines the core of a larger segment that will produce a significant spliced alignment. Our implementation is based on the initial identification of exactly matching 12mers by the suffix array method of Manber and Myers (1993). Matching 12mers are first maximally extended and then assembled into matching regions allowing for small insertions and deletions in both genomic and cDNA and longer gaps in the genomic DNA (possible introns). These regions in the genomic DNA are then extended by typically several hundred nucleotides to define the segment to which the sahmtD algorithm is

applied. Details will be presented elsewhere (W. Zhu and V. Brendel, in preparation). For our Web server we are periodically pre-processing the major publicly available plant EST collections. However, the pre-processing for user-specific EST collections can also be achieved in reasonable time. For example, pre-processing of 37 745 *Arabidopsis thaliana* ESTs from GenBank took 9 min 22 s on our server in single-user mode. Matching these ESTs onto a 107 kb contig (GenBank U89959) produced 89 spliced alignments in 8 min 29 s. For comparison, the sim4 program of Florea *et al.* (1998) took 18 min 35 s on the same data (note that this program produces all exact matches of length at least 12 and thus the output is much less specific than the GeneSeqer output; concerning sensitivity, see Figure 6 and Discussion).

#### Minimal intron length

The algorithm as given does not impose any restrictions on exon and intron lengths. Naturally occurring introns exceed a minimal length of about 55–60 bases (for plants, see Brendel *et al.*, 1998). To avoid solutions with unacceptably small intron assignments, the algorithm was modified to include a 'short intron penalty'. This penalty is levied upon intron to exon state transitions depending on whether or not the then closed intron exceeds the required minimum length. The implementation is straightforward: at the maximization step for  $I_m^n$ , a variable **intronstart**[ $n$ ][ $m$ ] is set to  $n$  (beginning of a new intron) or carried over from **intronstart**[ $n-1$ ][ $m$ ] (continuation of an existing intron; the index  $n$  of **intronstart** can again be replaced by  $n \bmod 2$ ). The weight is added during the  $E_m^n$  maximization if the current index  $n$  does not exceed the **intronstart** value by more than the defined minimal intron length. Because of the left to right directionality of the maximization algorithm, the modified procedure is not guaranteed to find the optimal score of all alignments satisfying the minimal intron length constraint. In practice, this poses no problem. The different donor site assignments are most likely well distinguished by the

(a)  
 Plus Strand HSPs:  
 Score = 916 (137.4 bits), Expect = 6.9e-53, Sum P(2) = 6.9e-53  
 Identities = 190/200 (95%), Positives = 190/200 (95%), Strand = Plus / Plus

Query: 2722 AGAGGGAAGCTCGACTGGAACGCAATAAGACACGGTCGTTGTTCTCGTGAATGCAGATCC 2781  
 AGAGGGAAGCTCGACTGGAACGCAATAAGACACGGTCGTTGTTCTCGTGAAT GCAGATCC  
 Sbjct: 172 AGAGGGAAGCTCGACTGGAACGCAATAAGACACGGTCGTTGTTCTCGTGAATGCAGATCC 231

Query: 2782 TCGAATACCAATGATGCTCTCAGAACATCACCTAGCTAGTAGTATCCTGTTGTTTCATTTG 2841  
 TCGAATACCAATGATGCTCTCAGAACATCACCTAGCTAGTAGTATCCTGTTGTTTCATTTG  
 Sbjct: 232 TCGAATACCAATGATGCTCTCAGAACATCACCTAGCTAGTAGTATCCTGTTGTTTCATTTG 291

Query: 2842 CAATGGCTGTGTTTGTATGATCTATCTAAGTAAACAAGTGAAGAGTGTGTTAATGTTA 2901  
 CAATGGCTGTGTTTGTATGATCTATCTAAGTAAACAAGTGAAGAGTGTGTTAATGTTA  
 Sbjct: 292 CAATGGCTGTGTTTGTATGATCTATCTAAGTAAACAAGTGAAGAGTGTGTTAATGTTA 351

Query: 2902 CTTTTTACCCCAATTGGT 2921  
 CTTTTTAC CCCC TTGG G  
 Sbjct: 352 CTTTTTACCCCC-FTGGNG 370

Score = 483 (72.5 bits), Expect = 6.9e-53, Sum P(2) = 6.9e-53  
 Identities = 97/98 (98%), Positives = 97/98 (98%), Strand = Plus / Plus

Query: 2186 GGGAAATGTCGACGAAAGCGCGCGCGCGGTACCCCTAGCCGGCTCGGATATCTGATT 2245  
 |||  
 Sbjct: 1 GGGAAATGTCGACGAAAGCGCGCGCGCGGTACCCCTAGCCGGCTCGGATATCTGATT 60

Query: 2246 CTCATGTTATCTTCAGTACTGCTCTCTCAAAATGT 2283  
 |||  
 Sbjct: 61 CTCATGTTATCTTCAGTACTGCTCTCTCAAAATGT 98

Score = 385 (57.8 bits), Expect = 1.7e-48, Sum P(2) = 1.7e-48  
 Identities = 103/128 (80%), Positives = 103/128 (80%), Strand = Plus / Plus

Query: 2514 TCCGTTTTCATTAGTTATGCCTCTTAGCTTGACCCCT-TGATT-TCTTATCAGGTCTTGA 2571  
 TC G T TC T A TT T C T TTA CTT A CT TG TT T T A A GT TTGA  
 Sbjct: 47 TCGGATATC-TGA-TTCTCCATGTTATCTTCAGTACTGCTCTCTCAAA-ATGTNTTGA 103

Query: 2572 AGAATTTGGATCAGACAAAGCTAAATGCCAGATCATTTTATGTTGTACAGGAATGCCAA 2631  
 AGAATTTGGATCAGACAAAGCTAAATGCCAGATCATTTTATGTTGTACAGGAATGCCAA  
 Sbjct: 104 AGAATTTGGATCAGACAAAGCTAAATGCCAGATCATTTTATGTTGTACAGGAATGCCAA 163

Query: 2632 GAAGAAAGAG 2641  
 GAAGAAAGAG  
 Sbjct: 164 GAAGAAAGAG 173

**Fig. 3.** Resolution of EST hits from a BLAST search. (a) The 3kb region 59001–62000 of the *Arabidopsis thaliana* contig U89959 ('Query') was subjected to a BLAST search against the *Arabidopsis* EST database using the Stanford Genome Center server (<http://genome-www2.stanford.edu/cgi-bin/AtDB/nph-blast2atdb>). Three hits are reported with the plus strand of the 371 base cDNA clone AA712564 ('Subject'). (b) SplicePredictor resolves the three BLAST hits into a single consistent spliced alignment consisting of three exons and two introns. The EST represents a full or almost full-length cDNA encoding a 71 amino acid polypeptide (start codon at 61191–61193, stop codon at 61769–61771). The scores for the predicted exons were calculated as described in the text. The predicted donor and acceptor sites are scored by *P*-value (Pd and Pa, respectively, Kleffe *et al.*, 1996) and by the similarity score (s) calculated for the proximal 50 exon bases. Alignment positions that align identical letters are indicated by vertical bars.

combination of *P*-value and alignment quality, and the intron length restriction would mainly serve to eliminate alignment paths that display within likely introns short stretches of relatively high sequence similarity that can be expected at random.

*Scoring the alignment*

The program scores each predicted exon separately by tallying up the output weights corresponding to the alignment of exon and cDNA. Only matches and penalties for gaps in the genomic DNA are counted. This value is normalized by the equivalent sum of weights assuming per-

(b)  
 Predicted gene structure:  
 Exon 1 61186 61281 ( 96 n); cDNA 1 96 ( 96 n); score: 0.990  
 Intron 1 61282 61564 ( 283 n); Pd: 0.020 (s: 0.98), Pa: 0.815 (s: 0.92)  
 Exon 2 61565 61641 ( 77 n); cDNA 97 173 ( 77 n); score: 0.948  
 Intron 2 61642 61723 ( 82 n); Pd: 0.067 (s: 0.94), Pa: 0.933 (s: 1.00)  
 Exon 3 61724 61922 ( 199 n); cDNA 174 371 ( 198 n); score: 0.945

CDS\_AA712564+:(61186..61281,61565..61641,61724..61922)

Alignment:  
 GGGAAATGTC GACGAAAGGC GCGCGCGCGG CGTACCCTAG CGCGGCTCGG ATATCTGATT 61245  
 |||  
 GGGAAATGTC GACGAAAGGC GCGCGCGCGG CGTACCCTAG CGCGGCTCGG ATATCTGATT 60

CTCCATGTTA TCTTCAGTAC TCTGCTTCTC TCAAATGTGA GTCATGCTCC TGATCTCAC 61305  
 |||  
 CTCCATGTTA TCTTCAGTAC TCTGCTTCTN TCAAAT..... 96

CTTTGTGATT GTTCTTCGGA GGATAGGATT TGACATGTTA TCTTCAGTAC TGTCAGTTC 61365  
 ..... 96

CATAACGAAT TAGCATTGAT TAGATCTCAT CTATTTTCATT ATGCTTCTTC AAGGTGATTA 61425  
 ..... 96

GATTAGTGGG TTGAATCCCA TGTCAGTATG TCGATTTAGG TCCCATCAAT TGATAACGTC 61485  
 ..... 96

GGGTTTGATT CCTGATTGTT TATGTGTTTC CGTTTTCATT AGTTATGCCT CTTAGCTTGA 61545  
 ..... 96

CCCCTTGATT TCTTATCAGG TCTTGAAGAA TTTGGATCAG ACAAGAGTAA ATGCCAGGAT 61605  
 |||  
 .....G TTTTGAAGAA TTTGGATCAG ACAAGAGTAA ATNCCAGGTT 137

CATTTTGATG TGTACAAGGA ATGCAAGAAG AAAGAGGTTG TTGTTGTGAA TGAATATTTA 61665  
 |||  
 CATTTTATG TGTACAAGGA ATGCAAGAAG AAAGAG..... 173

GGCTTTTGGC GTTTCACACT TCTTTGCTGC TTTACTCATG TGTATTTTGG TTTCTCAGAG 61725  
 |||  
 .....AG 175

GGAAAGCTCGA CTGGAACGCA ATAAAGACAG GTCCTTCTTC TCGTGAATGC AGATCCTCGA 61785  
 |||  
 GGAAGCTCGA CTGGAACGCA ATAAAGACAG GTCCTTCTTC TCGTGAATNC AGATCCTCGA 235

ATACCAATGA TGTCTCAGAA CATCACCTAG CTAGTAGTAT CCTGTTGTTT CATTTCGAAT 61845  
 |||  
 ATACCAATGA TGTCTCAGAA CATCACCTAG CTAGTAGTAT CCTGTTGTTT CATTTCGAAT 295

GGCTGTGTTT GTATGATCTA TCTAAGTAAA CAAGTGGAAA GTGTTGTGTTA ATGTTACTTT 61905  
 |||  
 GGCTGTGTTT GNATGANCTA TCTAAGTAAA CAAGTGGAAA GTTTTTNTNA ATGTTACTTT 355

TTACTCCCCA TTGGTGA 61922  
 |||  
 TTAC-CCCC TTGGNG 371

**Fig. 3.** cont.

fect matching to the genomic DNA. For ungapped alignments, this score is correlated with percentage identity. From our experience, the optimal alignment of unrelated ESTs to a genomic DNA rarely produces exon quality values above 0.4 for exons of lengths at least 60 nucleotides (data not shown).

For donor and acceptor sites, the program displays the *P*-values and evaluates the exon quality for the adjacent 50 exon bases. For non-cognate, but homologous ESTs, these values may indicate high conservation around the splice sites, even though the central parts of long exons may have diverged considerably.



**Table 1.** Predicted gene structure of two closely related *Arabidopsis thaliana* genes. The two potential genes located on chromosomes 1 (contig GenBank U89959) and 2 (contig GenBank AB008268) were predicted as described in the text. The ‘from/score’ and ‘to/score’ columns give the starting and ending positions of the exons or the splice site scores of the introns (see Brendel and Kleffe, 1998; 15\* is the optimal score). The ‘size’ column refers to the lengths of the exons and introns in number of nucleotides. The column ‘sim’ gives the similarity score comparing the corresponding exons from U89959 and AB008268 as derived from the spliced alignment (Figure 4). The similarity score for the first exon (shown in parenthesis) refers to the score for the 3’ end of the exons aligned as in Figure 4. Exons 4 and 5 as given by the coordinates in this Table were not aligned by the algorithm

		U89959			AB008268			
	#	from/score	to/score	size	from/score	to/score	size	sim
exon	1	62310	62646	337	20620	21025	406	(0.60)
intron	1	5*	10*	201	15*	15*	89	
exon	2	62848	63178	331	21115	21445	331	0.66
intron	2	12*	15*	92	15*	6*	81	
exon	3	63271	63391	121	21527	21644	118	0.60
intron	3	9*	3*	76	5*	15*	87	
exon	4	63468	63683	216	21732	21905	174	
intron	4	15*	15*	71	15*	10*	68	
exon	5	63755	63814	60	21974	22033	60	



**Fig. 5.** Alignment of the predicted protein sequences encoded by the exons from the *Arabidopsis thaliana* contigs GenBank U89959 and GenBank AB008268 as assigned in Table 1. The alignment was produced by the PPAT algorithm (V. Brendel, unpublished) which is an extension of the published SSPA algorithm (Karlin *et al.*, 1995). Introns are indicated by ‘=’. The alignment was scored with the BLOSUM62 amino acid substitution scoring matrix (Henikoff and Henikoff, 1992). Aligned residues are connected with a vertical bar if identical, by ‘+’ if positively scoring in the BLOSUM62 matrix, by ‘.’ if scoring 0, and by a blank if scoring negatively. Residues that could not be aligned in significantly scoring alignment blocks are given in lower case. It is seen that strong conservation extends over the C-terminal parts of exons 1, all of exons 2 and 3, the N-terminal parts of exons 4, and all of exons 5.

63580); exon boundaries in agreement with Table 1 are printed in bold face. This assignment includes non-consensus splice sites at 63181 (TA donor), 63403 (GA donor), and 63488 (CT acceptor). EST\_GENOME (version 4, obtained from ftp.sanger.ac.uk/pub/pmr) gives an alignment with only two introns, (62554..**62646,62848**..**63178,63271**..63370).

GenScan predicted the 5’ incomplete gene structure (62294..62481,62507..**62646,62848**..63138,63312..**63391,63432**..**63683,63755**..**63814**). This assignment includes the unrealistically small intron 63392 to 63431 of 40 nucleotides only. These comparisons suggest that our algorithm can usefully extend gene prediction in the presence of weak sequence similarity information. We



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